

## Original Paper

# Reducible Relationships between Generalized Syllogisms with the Quantifiers in Square{*at least half of the*}

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### Abstract

*The main conclusions of this paper are as follows: Theorem 1 proves the validity of the generalized syllogism SAI-3. Theorem 2 takes the syllogism SAI-3 as the fundamental axiom, and then deduces the other 14 non-trivial valid generalized syllogisms with the quantifiers in Square{at least half of the}. It means that there are reducible relationships between these 15 syllogisms, and that the above knowledge mining processes are consistent. This innovative achievement can promote in-depth research on generalized syllogistic and provide methodological inspiration for knowledge mining in artificial intelligence.*

### Keywords

*generalized syllogisms, validity, reducible relationships, knowledge mining*

## 1. Introduction

In natural language reasoning, in addition to common classical syllogism reasoning (Łukasiewicz, 1957), there is generalized syllogism reasoning (Hao, 2024a-b). A non-trivial generalized syllogism contains at least one non-trivial generalized quantifier (Peters & Westerståhl, 2006; Xu & Yu, 2024). There are only four classical quantifiers as follows: *some*, *no*, *not all*, and *all*, and they form Square{*some*} (Wang & Yuan, 2024). Non-trivial generalized quantifiers are non-classical ones, such as *at least half of the*, *fewer than half of the*, *at most half of the*, *most*. They form Square{*at least half of the*} (Ma & Cao, 2024).

Although there has been some works about generalized syllogisms (Moss, 2010; Endrullis & Moss, 2015; Cao & Li, 2024), there are still many interesting questions to be studied. This paper only studies the validity and reducibility of the generalized syllogisms involving the eight quantifiers from Square{*some*} and Square{*at least half of the*}.

## 2. Preliminaries

In this paper, let  $p$ ,  $t$ , and  $n$  be variables, and  $P$ ,  $T$ , and  $N$  be the sets that composed of  $p$ ,  $t$ ,  $n$ , respectively.  $D$  be the domain of variables. Let  $\lambda$ ,  $\omega$ ,  $\psi$ , and  $\delta$  be well-formed formulas (shorted as wff). ' $|P \cap N|$ ' represents the cardinality of the intersection of the set  $P$  and  $N$ , ' $\vdash \omega$ ' says that  $\omega$  is provable, and ' $\omega =_{\text{def}}$ '

$\delta'$  that  $\omega$  can be defined by  $\delta$ . The operators (such as  $\neg$ ,  $\rightarrow$ ,  $\wedge$ ,  $\leftrightarrow$ ) are basic ones in set theory (Halmos, 1974).

Let  $Q$  is a quantifier,  $\neg Q$ ,  $Q\neg$ , and  $\neg Q\neg$  are respectively its outer, inner, dual negative quantifier. These four quantifiers form a modern Square  $\{Q\} = \{Q, \neg Q, Q\neg, \neg Q\neg\}$  (Hao, 2024b). For example, Square  $\{some\} = \{some, no, not all, all\}$ , and Square  $\{at least half of the\} = \{at least half of the, fewer than half of the, at most half of the, most\}$ . The propositions containing these 8 quantifiers from Square  $\{some\}$  and Square  $\{at least half of the\}$  are respectively called Proposition  $I, E, O, A, S, F, H$ , and  $M$  (Ma & Cao, 2024). This paper only studies the non-trivial generalized syllogisms composed of these 8 propositions, and these syllogisms include at least one of the last four propositions. Thus, ' $at least half of the(t, n) \wedge all(t, p) \rightarrow some(p, n)$ ' is a third figure syllogism, which can be abbreviated as *SAI-3*.

Example 1:

Major premise: At least half of my flowers are roses.

Minor premise: All flowers are plants.

Conclusion: Some plants are roses.

Let  $t, n$ , and  $p$  be variables that represent a flower, a rose, and a plant, respectively, then this syllogism can be symbolized as ' $at least half of the(t, n) \wedge all(t, p) \rightarrow some(p, n)$ '.

Special note: It was assumed that the set composed of the subject variable for a categorical proposition containing one of the four quantifiers (*all, not all, most, and at most half of the*) is not an empty set.

### 3. Generalized Syllogism System with the quantifiers '*at least half of the*'

This system consists of the following parts:

#### 3.1 Primitive Symbols

- (1) variables:  $p, t, n$
- (2) quantifiers: *some, at least half of the*
- (3) operators:  $\neg, \rightarrow$
- (4) brackets:  $(, )$

#### 3.2 Formation Rules

- (1) If  $Q$  is a quantifier,  $p$  and  $n$  are variables, then  $Q(p, n)$  is a wff.
- (2) If  $\lambda$  and  $\omega$  are wffs, then so are  $\neg\lambda$  and  $\lambda \rightarrow \omega$ .
- (3) Only the sentences formed on the basis of the above two rules are wffs.

#### 3.3 Basic Axioms

A1: If  $\lambda$  is a valid proposition in classical logic, then  $\vdash \lambda$ .

A2:  $\vdash at least half of the(t, n) \wedge all(t, p) \rightarrow some(p, n)$  (i.e. the syllogism *SAI-3*).

#### 3.4 Rules of Deduction

Rule 1 (Antecedent strengthening): If  $\vdash (\lambda \wedge \omega \rightarrow \psi)$  and  $\vdash (\delta \rightarrow \lambda)$ , then  $\vdash (\delta \wedge \omega \rightarrow \psi)$ .

Rule 2 (subsequent weakening): If  $\vdash (\lambda \wedge \omega \rightarrow \psi)$  and  $\vdash (\psi \rightarrow \delta)$ , then  $\vdash (\lambda \wedge \omega \rightarrow \delta)$ .

Rule 3 (anti-syllogism): If  $\vdash (\lambda \wedge \omega \rightarrow \psi)$ , then  $\vdash (\neg\psi \wedge \lambda \rightarrow \neg\omega)$ .

### 3.5 Relevant Definitions

D1:  $(\omega \wedge \delta) =_{\text{def}} \neg(\omega \rightarrow \neg \delta)$ ;

D2:  $(\omega \leftrightarrow \delta) =_{\text{def}} (\omega \rightarrow \delta) \wedge (\delta \rightarrow \omega)$ ;

D3:  $(Q\neg)(p, n) =_{\text{def}} Q(p, D\neg n)$ ;

D4:  $(\neg Q)(p, n) =_{\text{def}}$  It is not the case that  $Q(p, n)$ ;

D5: *some*( $p, n$ ) is true iff  $P \cap N \neq \emptyset$  is true;

D6: *no*( $p, n$ ) is true iff  $P \cap N = \emptyset$  is true;

D7: *not all*( $p, n$ ) is true iff  $P \not\subseteq N$  is true;

D8: *all*( $p, n$ ) is true iff  $P \subseteq N$  is true;

D9: *at least half of the*( $p, n$ ) is true iff  $|P \cap N| \geq 0.5 |P|$  is true;

D10: *fewer than half of the*( $p, n$ ) is true iff  $|P \cap N| < 0.5 |P|$  is true;

D11: *at most half of the*( $p, n$ ) is true iff  $|P \cap N| \leq 0.5 |P|$ ;

D12: *most*( $p, n$ ) is true iff  $|P \cap N| > 0.5 |P|$  is true.

### 3.6 Relevant Facts

#### Fact 1 (inner negation):

(1.1)  $\vdash \text{all}(p, n) \leftrightarrow \text{no}\neg(p, n)$ ;

(1.2)  $\vdash \text{no}(p, n) \leftrightarrow \text{all}\neg(p, n)$ ;

(1.3)  $\vdash \text{some}(p, n) \leftrightarrow \text{not all}\neg(p, n)$ ;

(1.4)  $\vdash \text{not all}(p, n) \leftrightarrow \text{some}\neg(p, n)$ ;

(1.5)  $\vdash \text{most}(p, n) \leftrightarrow \text{fewer than half of the}\neg(p, n)$ ;

(1.6)  $\vdash \text{fewer than half of the}(p, n) \leftrightarrow \text{most}\neg(p, n)$ ;

(1.7)  $\vdash \text{at least half of the}(p, n) \leftrightarrow \text{at most half of the}\neg(p, n)$ ;

(1.8)  $\vdash \text{at most half of the}(p, n) \leftrightarrow \text{at least half of the}\neg(p, n)$ .

#### Fact 2 (outer negation):

(2.1)  $\vdash \neg \text{all}(p, n) \leftrightarrow \text{not all}(p, n)$ ;

(2.2)  $\vdash \neg \text{not all}(p, n) \leftrightarrow \text{all}(p, n)$ ;

(2.3)  $\vdash \neg \text{no}(p, n) \leftrightarrow \text{some}(p, n)$ ;

(2.4)  $\vdash \neg \text{some}(p, n) \leftrightarrow \text{no}(p, n)$ ;

(2.5)  $\vdash \neg \text{most}(p, n) \leftrightarrow \text{at most half of the}(p, n)$ ;

(2.6)  $\vdash \neg \text{at most half of the}(p, n) \leftrightarrow \text{most}(p, n)$ ;

(2.7)  $\vdash \neg \text{fewer than half of the}(p, n) \leftrightarrow \text{at least half of the}(p, n)$ ;

(2.8)  $\vdash \neg \text{at least half of the}(p, n) \leftrightarrow \text{fewer than half of the}(p, n)$ .

#### Fact 3 (symmetry):

(3.1)  $\vdash \text{some}(p, n) \leftrightarrow \text{some}(n, p)$ ;

(3.2)  $\vdash \text{no}(p, n) \leftrightarrow \text{no}(n, p)$ .

#### Fact 4 (subordination) :

(4.1)  $\vdash \text{all}(p, n) \rightarrow \text{some}(p, n)$ ;

- (4.2)  $\vdash no(p, n) \rightarrow not\ all(p, n);$
- (4.3)  $\vdash all(p, n) \rightarrow most(p, n);$
- (4.4)  $\vdash most(p, n) \rightarrow some(p, n);$
- (4.5)  $\vdash all(p, n) \rightarrow at\ least\ half\ of\ the(p, n);$
- (4.6)  $\vdash at\ least\ half\ of\ the(p, n) \rightarrow some(p, n);$
- (4.7)  $\vdash at\ least\ half\ of\ the(p, n) \rightarrow most(p, n);$
- (4.8)  $\vdash at\ most\ half\ of\ the(p, n) \rightarrow fewer\ than\ half\ of\ the(p, n);$
- (4.9)  $\vdash fewer\ than\ half\ of\ the(p, n) \rightarrow not\ all(p, n);$
- (4.10)  $\vdash at\ most\ half\ of\ the(p, n) \rightarrow not\ all(p, n);$
- (4.11)  $\vdash no(p, n) \rightarrow fewer\ than\ half\ of\ the(p, n);$
- (4.12)  $\vdash no(p, n) \rightarrow at\ most\ half\ of\ the(p, n).$

The above facts are the basis facts in generalized quantifier theory (Peters & Westerståhl, 2006).

#### 4. Knowledge Mining about Valid Generalized Syllogisms

The following Theorem 1 proves the validity of the generalized syllogism *SAI-3*. Taking the syllogism *SAI-3* as the fundamental axiom in Theorem 2, the other 14 valid generalized syllogisms can be derived by the above facts and definitions.

**Theorem 1 (*SAI-3*):** The generalized syllogism  $at\ least\ half\ of\ the(t, n) \wedge all(t, p) \rightarrow some(p, n)$  is valid.

Proof: Suppose that  $at\ least\ half\ of\ the(t, n)$  and  $all(t, p)$  are true, thus  $|T \cap N| \geq 0.5 |T|$  is true according to Definition D9, and  $T \subseteq P$  is true with the help of Definition D8. It follows that  $P \cap N \neq \emptyset$ . This can be proven using the method of contradiction. Suppose that  $P \cap N = \emptyset$ , thus  $T \cap N = \emptyset$  in line with  $T \subseteq P$ . This contradicts  $|T \cap N| \geq 0.5 |T|$ . Therefore,  $P \cap N \neq \emptyset$  holds. It can be concluded that  $some(p, n)$  is true in the light of Definition D5, just as desired.

**Theorem 2:** There are 14 valid generalized syllogisms inferred from *SAI-3*:

- (1)  $\vdash SAI-3 \rightarrow ASI-3$
- (2)  $\vdash SAI-3 \rightarrow HAO-3$
- (3)  $\vdash SAI-3 \rightarrow EAF-1$
- (4)  $\vdash SAI-3 \rightarrow EAF-1 \rightarrow EAF-2$
- (5)  $\vdash SAI-3 \rightarrow ESO-2$
- (6)  $\vdash SAI-3 \rightarrow ESO-2 \rightarrow ESO-1$
- (7)  $\vdash SAI-3 \rightarrow ASI-3 \rightarrow ESO-3$
- (8)  $\vdash SAI-3 \rightarrow ASI-3 \rightarrow ESO-3 \rightarrow ESO-4$
- (9)  $\vdash SAI-3 \rightarrow EAF-1 \rightarrow AAM-1$
- (10)  $\vdash SAI-3 \rightarrow EAF-1 \rightarrow EAF-2 \rightarrow AEF-2$
- (11)  $\vdash SAI-3 \rightarrow EAF-1 \rightarrow EAF-2 \rightarrow AEF-2 \rightarrow AEF-4$
- (12)  $\vdash SAI-3 \rightarrow ESO-2 \rightarrow AHO-2$
- (13)  $\vdash SAI-3 \rightarrow ESO-2 \rightarrow ESO-1 \rightarrow ASI-1$

(14)  $\vdash SAI-3 \rightarrow ESO-2 \rightarrow ESO-1 \rightarrow ASI-1 \rightarrow SAI-4$

Proof:

- [1]  $\vdash \text{at least half of the}(t, n) \wedge \text{all}(t, p) \rightarrow \text{some}(p, n)$  (i.e. *SAI-3*, basic axiom A2)
- [2]  $\vdash \text{at least half of the}(t, n) \wedge \text{all}(t, p) \rightarrow \text{some}(n, p)$  (i.e. *ASI-3*, by [1] and Fact (3.1))
- [3]  $\vdash \text{at most half of the}(t, n) \wedge \text{all}(t, p) \rightarrow \text{not all}(p, n)$  (by [1], Fact (1.7) and (1.3))
- [4]  $\vdash \text{at most half of the}(t, D-n) \wedge \text{all}(t, p) \rightarrow \text{not all}(p, D-n)$  (i.e. *HAO-3*, by [1] and Definition D3)
- [5]  $\vdash \neg \text{some}(p, n) \wedge \text{all}(t, p) \rightarrow \neg \text{at least half of the}(t, n)$  (by [1] and Rule 3)
- [6]  $\vdash \text{no}(p, n) \wedge \text{all}(t, p) \rightarrow \text{fewer than half of the}(t, n)$  (i.e. *EAF-1*, by [5], Fact (2.4) and (2.8))
- [7]  $\vdash \text{no}(n, p) \wedge \text{all}(t, p) \rightarrow \text{fewer than half of the}(t, n)$  (i.e. *EAF-2*, by [6] and Fact (3.2))
- [8]  $\vdash \neg \text{some}(p, n) \wedge \text{at least half of the}(t, n) \rightarrow \neg \text{all}(t, p)$  (by [1] and Rule 3)
- [9]  $\vdash \text{no}(p, n) \wedge \text{at least half of the}(t, n) \rightarrow \text{not all}(t, p)$  (i.e. *ESO-2*, by [8], Fact (2.4) and (2.1))
- [10]  $\vdash \text{no}(n, p) \wedge \text{at least half of the}(t, n) \rightarrow \text{not all}(t, p)$  (i.e. *ESO-1*, by [9] and Fact (3.2))
- [11]  $\vdash \text{at least half of the}(t, n) \wedge \text{no}(t, p) \rightarrow \text{not all}(n, p)$  (by [2], Fact (1.1) and (1.3))
- [12]  $\vdash \text{at least half of the}(t, n) \wedge \text{no}(t, D-p) \rightarrow \text{not all}(n, D-p)$  (i.e. *ESO-3*, by [11] and Definition D3)
- [13]  $\vdash \text{at least half of the}(t, n) \wedge \text{no}(D-p, t) \rightarrow \text{not all}(n, D-p)$  (i.e. *ESO-4*, by [12] and Fact (3.2))
- [14]  $\vdash \text{all}(p, n) \wedge \text{all}(t, p) \rightarrow \text{most}(t, n)$  (by [6], Fact (1.2) and (1.6))
- [15]  $\vdash \text{all}(p, D-n) \wedge \text{all}(t, p) \rightarrow \text{most}(t, D-n)$  (i.e. *AAM-1*, by [14] and Definition D3)
- [16]  $\vdash \text{all}(n, p) \wedge \text{no}(t, p) \rightarrow \text{fewer than half of the}(t, n)$  (by [7], Fact (1.1) and (1.2))
- [17]  $\vdash \text{all}(n, D-p) \wedge \text{no}(t, D-p) \rightarrow \text{fewer than half of the}(t, n)$  (i.e. *AEF-2*, by [16] and Definition D3)
- [18]  $\vdash \text{all}(n, D-p) \wedge \text{no}(D-p, t) \rightarrow \text{fewer than half of the}(t, n)$  (i.e. *AEF-4*, by [17] and Fact (3.2))
- [19]  $\vdash \text{all}(p, n) \wedge \text{at most half of the}(t, n) \rightarrow \text{not all}(t, p)$  (by [9], Fact (1.2) and (1.7))
- [20]  $\vdash \text{all}(p, D-n) \wedge \text{at most half of the}(t, D-n) \rightarrow \text{not all}(t, p)$  (i.e. *AHO-2*, by [19] and Definition D3)
- [21]  $\vdash \text{all}(n, p) \wedge \text{at least half of the}(t, n) \rightarrow \text{some}(t, p)$  (by [10], Fact (1.2) and (1.4))
- [22]  $\vdash \text{all}(n, D-p) \wedge \text{at least half of the}(t, n) \rightarrow \text{some}(t, D-p)$  (i.e. *ASI-1*, by [21] and Definition D3)
- [23]  $\vdash \text{all}(n, D-p) \wedge \text{at least half of the}(t, n) \rightarrow \text{some}(D-p, t)$  (i.e. *SAI-4*, by [22] and Fact (3.1))

The above proof processes are logical deduction ones. Therefore, the above knowledge mining processes are consistent. Theorem 2 demonstrates that the validity of the above 14 non-trivial generalized syllogisms can be derived from the validity of the syllogism *SAI-3*. In other words, there are reducible relationships between these 15 syllogisms with the quantifiers in Square{*at least half of the*}.

## 5. Conclusion

The main conclusions of this paper are as follows: Theorem 1 proves the validity of the generalized syllogism *SAI-3*. Theorem 2 takes the syllogism *SAI-3* as the fundamental axiom, and then deduces the other 14 non-trivial valid generalized syllogisms with the quantifiers in Square{*at least half of the*}. It means that there are reducible relationships between these 15 syllogisms, and that the above knowledge mining processes are consistent.

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